

Introduction to Data Structure



Chapter 1: Introduction

A decorative graphic at the bottom of the slide consists of a horizontal bar with a color gradient from magenta on the left to dark purple on the right. The right end of the bar is shaped into a large arrow pointing to the right. Inside this arrow, there is a smaller, nested arrow shape, also pointing right, which is white with a magenta outline.

Abstract Data Types:

ADT

- ✓ A set of data values and associated operations that are precisely specified independent of any particular implementation.
- ✓ Example: stack, queue, priority queue.

DataStructures

- ✓ The term data structure refers to a scheme for organizing related pieces of Information.
- ✓ The basic types of data structures include: files, lists, arrays, records, trees, tables.
- ✓ A data structure is the concrete implementation of that type, specifying how much memory is required and, crucially, how fast the execution of each operation will be

Timing

- ✓ Every time we run the program we need to estimate how long a program will run since we are going to have different input values so the running time will vary.
- ✓ The worst case running time represents the maximum running time possible for all input values.
- ✓ We call the worst case timing big Oh written $O(n)$. The n represents the worst case execution time units.

Complexities

- ✓ Simple programming statement

- ✓ Example:

```
k++;
```

- ✓ Complexity : $O(1)$
- ✓ Simple programming statements are considered 1 time unit.

Complexities

- ✓ Linear *for* loops

- ✓ Example:

```
k = 0;
for(i = 0; i < n; i++)
    k++;
```

- ✓ Complexity : $O(n)$
- ✓ *for* loops are considered n time units because they will repeat a programming statement n times.
- ✓ The term linear means the for loop increments or decrements by 1

Complexities

- ✓ Non Linear loops

- ✓ Example:

```
k=0;
for(i=n; i>0; i=i/2)
    k++;
```

- ✓ Complexity : $O(\log n)$
- ✓ For every iteration of the loop counter i will divide by 2.
- ✓ If i starts is at 16 then then successive i 's would be 16, 8, 4, 2, 1. The final value of k would be 4. Non linear loops are logarithmic.
- ✓ The timing here is definitely $\log_2 n$ because $2^4 = 16$. Can also works for multiplication.

Complexities

- ✓ Nested *for* loops

- ✓ Example:

```
k=0;
for(i=0; i<n; i++)
  for(j=0; j<n; j++)
    k++;
```

- ✓ Complexity : $O(n^2)$

- ✓ Nested for loops are considered n^2 time units because they represent a loop executing inside another loop.

Complexities

- ✓ Sequential *for* loops

- ✓ Example:

```
k=0;
for(i=0; i<n; i++)
    k++;
k=0;
for(j=0; j<n; j++)
    k++;
```

- ✓ Complexity : $O(n)$
- ✓ Sequential for loops are not related and loop independently of each other.

Complexities

- ✓ Loops with non-linear inner loops

- ✓ Example:

```
k=0;
for(i=0;i<n;i++)
    for(j=i; j>0; j=j/2)
        k++;
```

- ✓ Complexity : $O(n \log n)$
- ✓ The outer loop is $O(n)$ since it increments linear.
- ✓ The inner loop is $O(n \log n)$ and is non-linear because decrements by dividing by 2.
- ✓ The final worst case timing is: $O(n) * O(\log n) = O(n \log n)$

Complexities

- ✓ Inner loop incrementer initialized to outer loop incrementer

- ✓ Example:

```
k=0;
for(i=0;i<n;i++)
    for(j=i;j<n;j++)
        k++;
```

- ✓ Complexity : $O(n^2)$
- ✓ In this situation we calculate the worst case timing using both loops.
- ✓ For every i loop and for start of the inner loop j will be n-1 , n-2, n-3.

Complexities

- ✓ Power Loops

- ✓ Example:

```
k=0;
for(i=1;i<=n;i*i*2)
  for(j=1;j<=i;j++)
    k++;
```

- ✓ Complexity : $O(2^n)$

- ✓ To calculate worst case timing we need to combine the results of both loops.
- ✓ For every iteration of the loop counter i will multiply by 2.
- ✓ The values for j will be 1, 2, 4, 8, 16 and k will be the sum of these numbers 31 which is $2^n - 1$.

Complexities

- ✓ If-else constructs
- ✓ Example:
- ✓ Complexity : $O(n^2)$

```
/* O(n) */  
if (x == 5)  
{  
    k=0;  
    for(i=0; i<n; i++)  
        k++;  
}  
/* O(n2) */  
else  
{  
    k=0;  
    for(i=0;i<n;i++)  
        for(j=i; j>0; j=j/2)  
            k++;  
}
```

Complexities

✓ Recursive

✓ Example:

```
int f(int n)
{
    if(n == 0)
        return 0;
    else
        return f(n-1) + n
}
```

✓ Complexity : $O(n)$

Stages: Program Design

- ✓ Identify the data structures.
- ✓ Operations: Algorithms
- ✓ Efficiency(Complexity)
- ✓ Implementation
 - ✓ What goes into header file?
 - ✓ What goes into C program?
 - ✓ What are libraries? Why do we need them?
 - ✓ How to create libraries.